

Bounds for D -Radius and D -diameter of a graphs

D. Reddy Babu and P. L. N. Varma

Communicated by Ayman Badawi

MSC2010 classification: primary 05C12.

keywords: D -distance, D -radius, D -diameter, maximum degree, minimum degree

Abstract. In this article, we give bounds for D -radius and D -diameter of a graph. Also, as D -radius, D -diameter is greater than radius and diameter we give bounds for their difference.

1 Introduction

The concept of distance is one of the important parameters in a graph. In a graph, geodesic distance is well known. In an article [2], the first two authors introduced the concept of D -distance between vertices by considering degree of vertices present in the path also. Using this distance we can define the eccentricity of a vertex and radius, diameter of any graph.

In this article we obtained bounds for D -radius and D -diameter in terms of minimal degree, maximal degree of vertices and number of vertices . Also we proved examples to show that the above bounds are attained.

Further we give an upper bound for the difference of D -radius and radius, D -diameter and diameter and show that in case of cyclic graph this bounds attained.

For any connected graph we use Δ to denote the maximal degree of vertices and δ to denote minimal degree of vertex. Further, we denote radius and diameter by $r(G)$ and $d(G)$ of a graph and $r^D(G)$, $d^D(G)$ to denote the D -radius and D -diameter of a graph.

2 Main results

Theorem 2.1. *If any graph G with n vertices and $\Delta(G) \geq 2$, then $2\delta + 1 \leq r^D(G) \leq d^D(G) \leq n(\Delta + 1) - 4$.*

Proof. Let u and v be any two vertices of G such that $d^D(u, v) = d^D(G)$ and $(u = u_1, u_2, u_3, \dots, u_m = v)$ be the corresponding path, where $m \leq n$. Then $d^D(u, v) = \sum deg(u) + m - 1 = m\Delta + (m - 1) = m(\Delta + 1) - 1 \leq n(\Delta + 1) - 1 \Rightarrow d^D(u, v) \leq n(\Delta + 1) - 1$. Now we claim that there is no graph G having $d^D(G) = n(\Delta + 1) - 1$. Suppose G is a graph having $d^D(G) = n(\Delta + 1) - 1$. This equality is obtained by taking the length of the D -path m as n and all vertices in the corresponding path are of degree Δ . Hence G must be a Δ - regular graph and $(u_1, u_2, u_3, \dots, u_n)$ is longest D -path in G . Now as $deg(u_i) = \Delta$, $\Delta \geq 2$, the vertex u_1 must adjacent to a vertex u_k , where $k \geq 2$. Then $(u_1, u_k, u_{k+1}, \dots, u_n)$ is a D -path between u_1 and u_n , of D -length less than $d^D(u, v)$. This is contradiction and so no graph can have the D -diameter $n(\Delta + 1) - 1$. Therefore $d^D(G) \leq n(\Delta + 1)$ Now let $x, y \in V(G)$ such that $d^D(x, y) = r^D(G)$ and $(x = v_1, v_2, v_3, \dots, v_l = y)$ be the corresponding path with $l \geq 2$. Then $d^D(x, y) = \sum deg(v_i) + l - 1 \geq l\delta + l - 1 \geq r(\delta + 1) - 1 = 2(\delta + 1) - 1 = 2\delta + 1$.

Therefore $2\delta + 1 \leq r^D(G) \leq d^D(G) \leq n(\Delta + 1) - 4$. □

Theorem 2.2. *For any graph G with $\Delta(G) \geq 2$, $d^D(G) = n(\Delta + 1) - 4$, if and only if $G = C_3$ or C_4 .*

Proof. Let $d^D(G) = n(\Delta + 1) - 4$ and p be a D - path in G of length $d^D(G)$.

Claim: the path p can exclude one vertex of G not belongs to p . Then $d^D(G) \leq (n - 2)(\Delta + 1) - (n - 3) = n(\Delta + 1) - 2(\Delta + 1) - n + 3 = n(\Delta + 1) - 2\Delta + 1 - n = n(\Delta + 1) - 4 - 2\Delta + 5 - n$ if $-2\Delta + 5 - n \geq 0$ then $2\Delta - 5 + n \leq 0 \Rightarrow 2\Delta + n \leq 5$, this is a contradiction and as $-2\Delta - n + 5 < 0$. Therefore $d^D(G) < n(\Delta + 1) - 4$ which is contradiction to our assumption and

so p can exclude at most one vertex of G . Now in order to prove the assumption, it is enough to prove that $p(G)$ does not contain all the vertices of G . Suppose $p(G)$ contain all vertices of G . If there is a vertex w such that $\deg(w) \geq 3$, then $p(G)$ will not contain all vertices of neighbor of w and so $\Delta(G) \leq 2$. Since $\Delta(G) \geq 2$, we have $\Delta(G) = 2$. Further if $\deg(u) = 1$ for some u , then $d^D(G) < n(\Delta + 1) - 4$ and so G must be cycle. since D -path between any two vertices of a cycle C_n will not include all the vertices of C_n and $p(G = C_n)$ contains all the vertices of $G = C_n$. we get a contradiction and hence the claim follows. Now $d^D(G) \leq (n - 1)(\Delta + 1) + n - 2 = (n - 1)(\Delta + 1) - 1$. By hypothesis $d^D(G) = n(\Delta + 1) - 4 \leq (n - 1)(\Delta + 1) - 1 \Rightarrow \Delta = 2$. Therefore $d^D(G) = n(\Delta + 1) - 4$ by theorem (3.2) in [6]. $n = 3$ or 4 . Thus G is either C_3 or C_4 . Converse is obvious. \square

Theorem 2.3. For any graph G , then $r^D(G) = 2\delta + 1$, if and only if $G \approx K_n$.

Proof. Let $r^D(G) = 2\delta + 1$ and $e^D(u) = 2\delta + 1, \forall u \in V(G)$, let $d^D(u, v) = 2\delta + 1$.

Claim: u and v are adjacent vertices. Suppose u and v are not adjacent is an internal vertex w in the $u-v$ D -path. Now $d^D(u, v) = d(u, v) + \deg(u) + \deg(v) + \deg(w) = 2 + 3\delta$. $2\delta + 1 = 3\delta + 1$ which is contradiction. Therefore there is no internal vertex in $u - v$ path and hence u and v are adjacent vertices. Therefore G is a complete graph. converse is obvious. \square

Next, we prove bounds for difference of radius and diameters.

Theorem 2.4. For any graph G , $r^D(G) - r(G) \leq \Delta(r(G) + 1)$.

Proof. Let v be a vertex of G such that $e(v) = r(G)$. Clearly $r^D(G) \leq e^D(v)$. Now u be farthest vertex from v with respect to D -distance and $d(u, v) = m$ then $e^D(u) \leq (m + 1)\Delta + m = m(\Delta + 1) + \Delta < e(v)(\Delta + 1) + \Delta$. $r^D(G) - r(G) \leq e(v)(\Delta + 1) + \Delta e(v) = \Delta(e(v) + 1) = \Delta(r(G) + 1)$. Therefore $r^D(G) - r(G) \leq \Delta(r(G) + 1)$. \square

Theorem 2.5. For any graph G , $d^D(G) - d(G) \leq \Delta(d(G) + 1)$.

Proof. Let v be a vertex of G such that $e^D(v) = d^D(G)$ and u be the farthest vertex v with respect to D -distance. let $d(u, v) = m$ then $d^D(v) = e^D(v) \leq (m + 1)\Delta + m = m(\Delta + 1) + \Delta = m(\Delta + 1) + \Delta$
 $d^D(G) \leq d(G)(\Delta + 1) + \Delta - d(G) = \Delta(d(G) + 1)$. And hence $d^D(G) - d(G) \leq \Delta(d(G) + 1)$. \square

Theorem 2.6. For any graph G , the following are equivalent.

- (i) $r^D(G) - r(G) = \Delta(r(G) + 1)$
- (ii) $d^D(G) - d(G) = \Delta(d(G) + 1)$
- (iii) $G = C_n$.

Proof. (i) \Rightarrow (ii)

Let us assume that $r^D(G) - r(G) = \Delta(r(G) + 1)$.

$$r(G) = \lceil \frac{n-1}{2} \rceil.$$

Suppose $r(G) < \lceil \frac{n-1}{2} \rceil$, let u be the vertex of G such that $e(u) = r(G)$ and v be a farthest vertex of u with respect to D -distance then $r^D(G) \leq e^D(u) \leq (d(u, v) + 1)\Delta + d(u, v) = d(u, v)(\Delta + 1) + \Delta \leq e(u)(\Delta + 1) + \Delta = r(G)(\Delta + 1) + \Delta$
 $r^D(G) - r(G) \leq r(G)(\Delta + 1) + \Delta - r(G) = \Delta(r(G) + 1) < \Delta(\lceil \frac{n-1}{2} \rceil + 1)$, which is contradiction and hence the assumption follows. Therefore G must be either a path or cycle. Suppose G is path on n vertices.

Case(1) n is even $r^D(G) = \frac{3n+2}{2}$ and $r(G) = \frac{n}{2}$

$$r^D(G) - r(G) = \frac{3n+2}{2} - \frac{n}{2} = n + 1$$

$(n + 1) < \Delta(r(G) + 1)$ which is contradiction.

Case(2) n is odd $r^D(G) = \frac{3n-1}{2}$ and $r(G) = \frac{n-1}{2}$

$$r^D(G) - r(G) = \frac{3n-1}{2} - \frac{n-1}{2} = n$$

$n < \Delta(r(G) + 1)$ which is also contradiction. Therefore G must be a cycle.

(iii) \Rightarrow (i)

If n is even $r^D(C_n) = \frac{3n+4}{2}$ and $r(C_n) = \frac{n}{2}$

$$r^D(C_n) - r(C_n) = \frac{3n+4}{2} - \frac{n}{2} = n + 2$$

$$\Delta(r(G) + 1) = 2(\frac{n}{2} + 1) = n + 2.$$

when n is odd $r^D(C_n) = \frac{3n+1}{2}$ and $r(C_n) = \frac{n-1}{2}$

$$r^D(C_n) - r(C_n) = \frac{3n+1}{2} - \frac{n-1}{2} = n + 1$$

$$\Delta(r(G) + 1) = 2\left(\frac{n-1}{2} + 1\right) = n + 1.$$

(iii) \Rightarrow (ii)

Let G be a cycle on n vertices then, $d^D(C_n) = \frac{3n+4}{2}$ and $d(C_n) = \frac{n}{2}$

$$d^D(C_n) - d(C_n) = \frac{3n+4}{2} - \frac{n}{2} = n + 2, \text{ if } n \text{ is even}$$

$$\Delta(d(G) + 1) = 2\left(\frac{n}{2} + 1\right) = n + 2. \text{ And}$$

$d^D(C_n) = \frac{3n+1}{2}$ and $d(C_n) = \frac{n-1}{2}$, if n is odd

$$d^D(C_n) - d(C_n) = \frac{3n+1}{2} - \frac{n-1}{2} = n + 1$$

$$\Delta(d(G) + 1) = 2\left(\frac{n-1}{2} + 1\right) = n + 1. \text{ Hence (ii) follows.}$$

(ii) \Rightarrow (iii)

We have to prove that $G = C_n$.

Claim: $deg(v_i) = \Delta(G) \forall v_i \in p$. p be the D -path suppose $deg(v_i) < \Delta$, for some vertex v_i in D -path p . Then $d^D(G) = \sum deg(v_i) + d(u, v) < d(u, v)\Delta + d(u, v) = d(u, v)(\Delta + 1) < d(G)(\Delta + 1)$, where $d(G)$ is the diameter of G . Therefore $d^D(G) - d(G) < d(G)(\Delta + 1) - d(G) = \Delta d(G)$ and so $\Delta(d(G) + 1) < \Delta d(G)$. This is contradiction as $n > d(G)$. Hence $deg(v_i) = \Delta$, for every v_i in p . Therefore $G = C_n$. \square

References

- [1] F. Buckley and F. Harary, *Distance in Graph*, Addison-Wesley, Longman, 1990.
- [2] D. Reddy Babu and P.L.N. Varma, *D-distance in graphs*, Golden Research Thoughts, 2 (9)(2013), 53-58.
- [3] D. Reddy Babu and P.L.N. Varma, *Average D-distance between edges of a graph*, Indian journal of science and technology, 8 (2)(2015), 152-156.
- [4] D. Reddy Babu and P.L.N. Varma, *Vertex-to-edge centers w.r.t D-distance*, Italian journal of pure and applied mathematics, 35 (2015), 101-108.
- [5] D. Reddy Babu and P.L.N. Varma, *A note on radius and diameter of graph w.r.t. D-distance*, International journal of chemical Science, 14(7)(2016), 1725-1729.
- [6] D. Reddy Babu and P.L.N. Varma, *D-radius and D-diameter of some families of graphs*, International journal of chemical science, 15(1) (2017), 1-9.

Author information

D. Reddy Babu, Department of A. S and H, Tirumala Engineering college, Jonnalagadda-522601, Narasaraopet, Guntur, India.

E-mail: reddybabu17@gmail.com

P. L. N. Varma, Division of mathematics, Department of Science & Humanities, V. F. S. T. R, Vadlamudi- 522 237, Guntur, India.

E-mail: plnvarma@gmail.com

Received: December 17, 2020

Accepted: May 16, 2021